Measurable Brooks's Theorem for Directed Graphs

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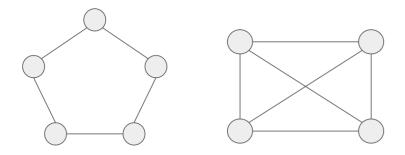
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Theorem (Brooks's theorem)

Suppose each vertex of G has degree at most $d \ge 2$. If d = 2, suppose G has no odd cycles; if $d \ge 3$, suppose G does not contain the complete graph on d + 1 vertices. Then $\chi(G) \le d$.

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Fact

[Kechris-Solecki–Todorčević, 1999] There is a Borel graph G for which $\chi(G) = 2$ but $\chi_B(G)$ is uncountable.

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Suppose each vertex of G has degree at most $d \ge 2$ and G contains no odd cycles or complete graphs. Then is $\chi_{\mathbf{B}}(G) \le d$?

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In fact, the set of acyclic *d*-regular Borel graphs having Borel chromatic number at most *d* is Σ_2^1 -complete [Brandt-Chang-Grebík-Grunau-Rozhoň-Vidnyánszky, 2024].

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Theorem (*µ*-measurable Brooks's theorem)

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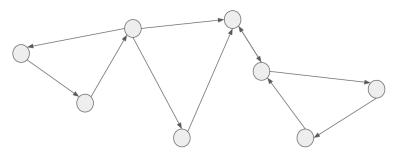
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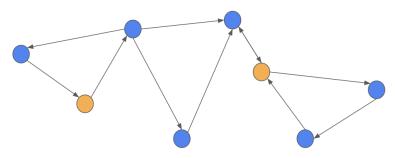


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[Harutyunyan-Mohar, 2011] Suppose the maximum degree of each vertex in D is at most $d \ge 2$. If d = 2, suppose D has no symmetrizations of odd cycles; if $d \ge 3$, suppose D does not contain the symmetrization of the complete graph on d + 1 vertices. Then $\overrightarrow{\chi}(D) \le d$.

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The directed Schreier graph of this action with generators $\gamma_0, \gamma_1, \gamma_2$ has no Borel 3-dicoloring.

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- Olicolor the "hard" components, the infinite Gallai trees, by discarding a μ-null set.

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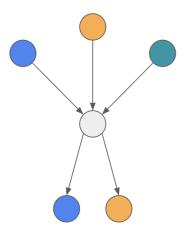
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- A connected digraph is a Gallai tree if all its blocks are bad.

Gallai Trees

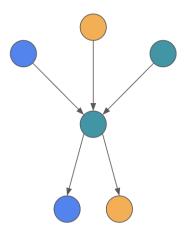
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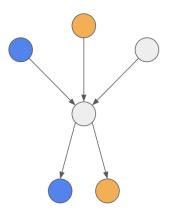


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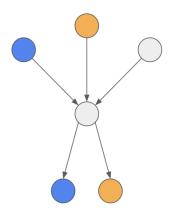


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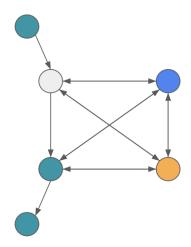
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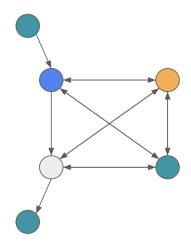


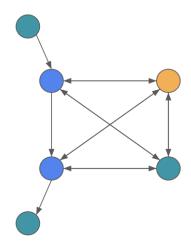
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Select the reserved vertices from an independent set of vertices with minimum degree less than d.







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Select the reserved good blocks so that, for distinct selected blocks S, T, no neighbor of S is also a neighbor of T.

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It may be instructive to consider the μ -measurable or Baire-measurable setting.

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- What can we say about LOCAL algorithms for dicoloring? Are there implications between descriptive digraph combinatorics and LOCAL dicoloring algorithms?

Thank you!